

RAMSEY GOODNESS OF BOUNDED DEGREE TREES

Igor Balla

Department of Mathematics, ETH Zürich

Joint work with: Alexey Pokrovskiy, Benny Sudakov

Ramsey number: $R(G, H)$ is the minimum N such that any red-blue coloring of K_N contains either a red copy of G or a blue copy of H .

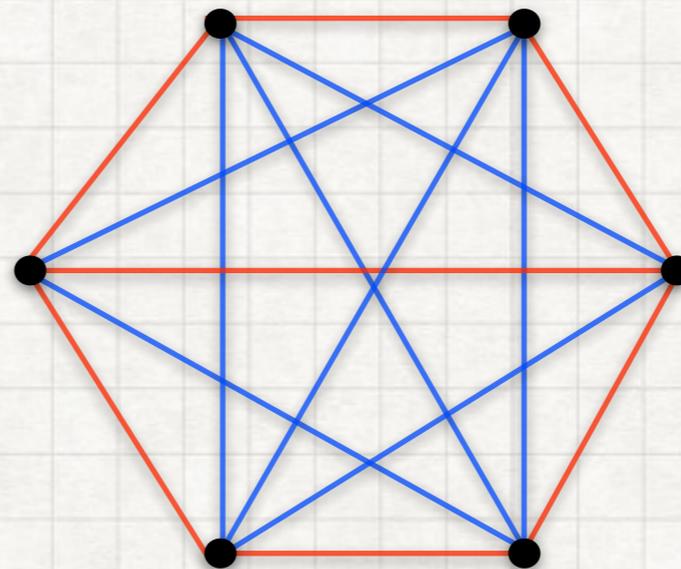
Complete graph on N vertices



Ramsey number: $R(G, H)$ is the minimum N such that any red-blue coloring of K_N contains either a red copy of G or a blue copy of H .

Complete graph on N vertices

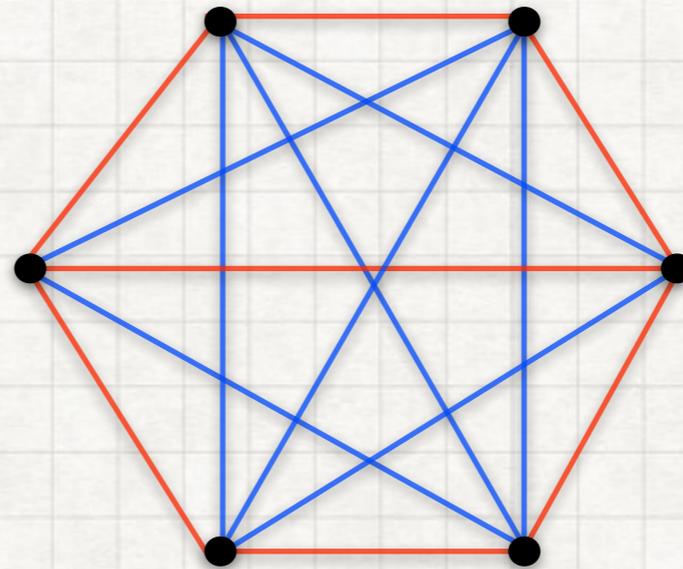
$$R(K_3, K_3) = 6$$



Ramsey number: $R(G, H)$ is the minimum N such that any red-blue coloring of K_N contains either a red copy of G or a blue copy of H .

Complete graph on N vertices

$$R(K_3, K_3) = 6$$

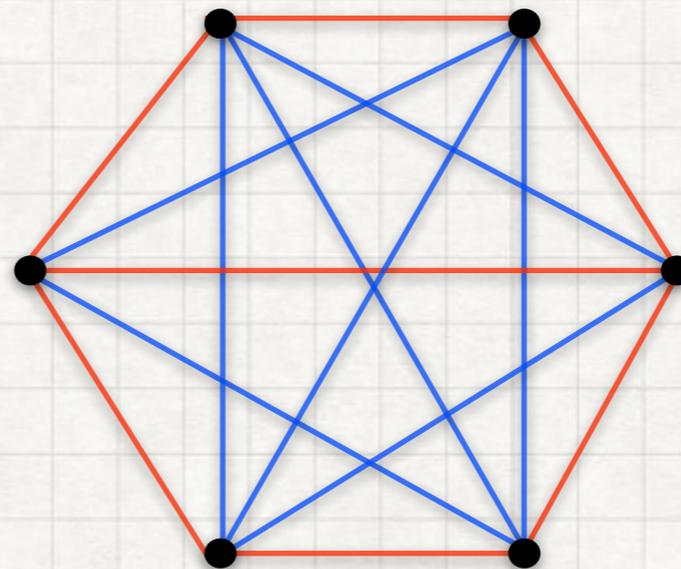


Theorem (Ramsey 1930): $R(K_n, K_n)$ is well defined.

Ramsey number: $R(G, H)$ is the minimum N such that any red-blue coloring of K_N contains either a red copy of G or a blue copy of H .

Complete graph on N vertices

$$R(K_3, K_3) = 6$$



Theorem (Ramsey 1930): $R(K_n, K_n)$ is well defined.

Theorem (Erdős 1947; Erdős and Szekeres 1935):

$$(\sqrt{2})^n \leq R(K_n, K_n) \leq 4^n$$

Theorem (Erdős 1947): $R(P_n, K_m) = (n - 1)(m - 1) + 1$

Path with n vertices

Theorem (Erdős 1947): $R(P_n, K_m) = (n - 1)(m - 1) + 1$

Path with n vertices

Theorem (Chvatal 1977): $R(T_n, K_m) = (n - 1)(m - 1) + 1$

Tree with n vertices

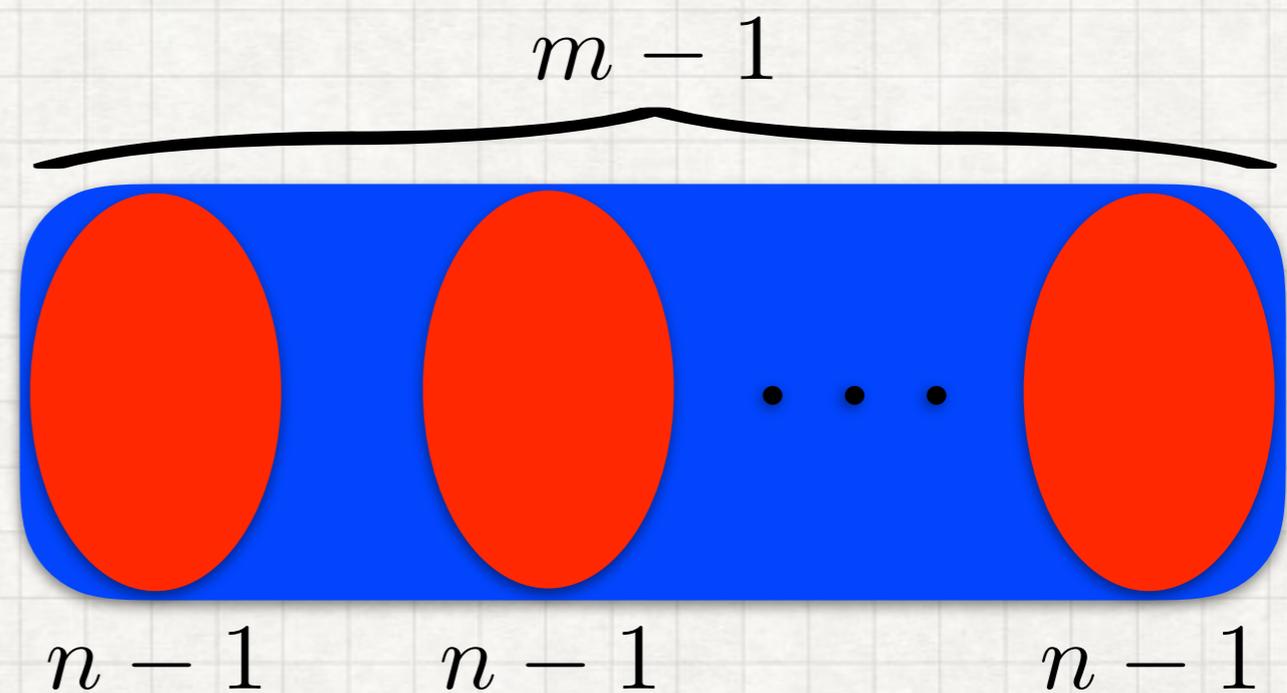
Theorem (Erdős 1947): $R(P_n, K_m) = (n - 1)(m - 1) + 1$

Path with n vertices

Theorem (Chvatal 1977): $R(T_n, K_m) = (n - 1)(m - 1) + 1$

Tree with n vertices

Lower bound construction:



Ramsey Goodness

A large grid of graph paper, consisting of 20 columns and 20 rows of small squares, intended for calculations or drawing.

Ramsey Goodness

- $\chi(H)$ = smallest number of colors in a proper coloring of H
- $\sigma(H)$ = minimum size of a color class in a $\chi(H)$ -coloring of H

Ramsey Goodness

- $\chi(H)$ = smallest number of colors in a proper coloring of H
- $\sigma(H)$ = minimum size of a color class in a $\chi(H)$ -coloring of H

Theorem (Burr 1981): For any connected G with $|G| \geq \sigma(H)$

$$R(G, H) \geq (|G| - 1)(\chi(H) - 1) + \sigma(H)$$

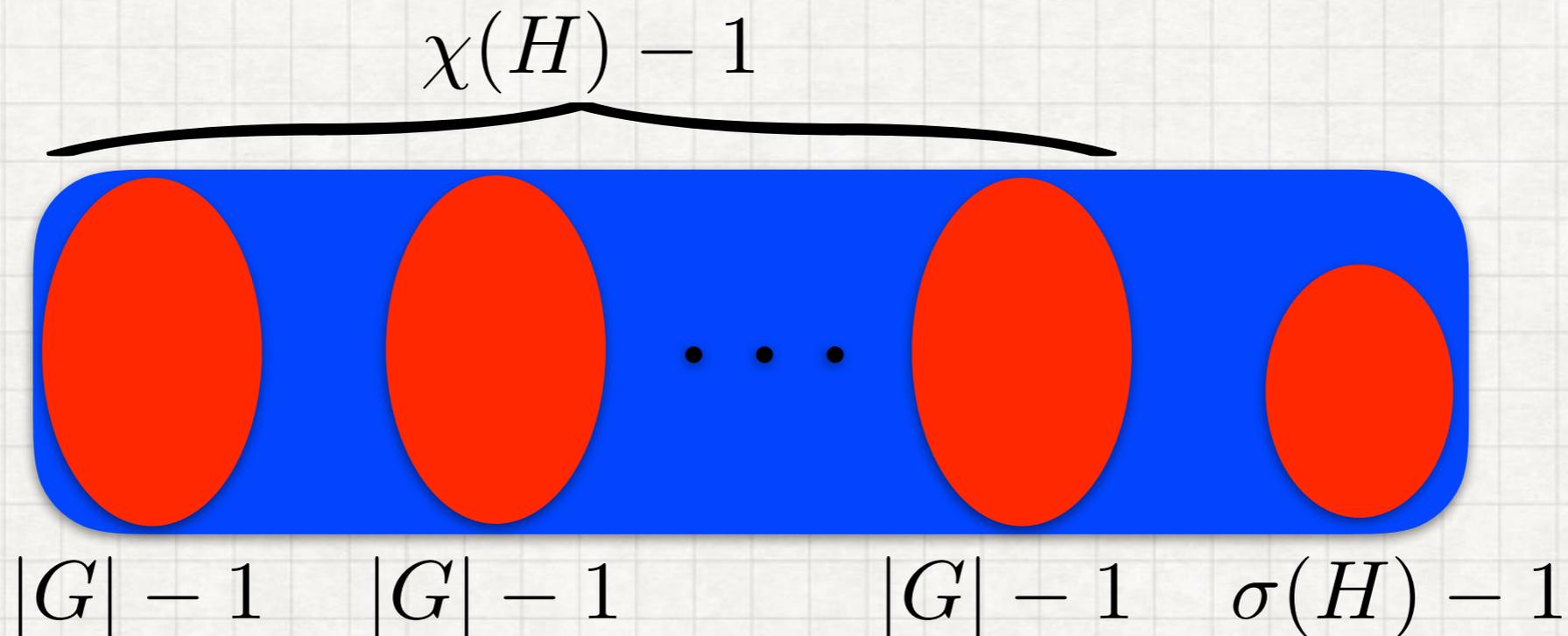
Ramsey Goodness

- $\chi(H)$ = smallest number of colors in a proper coloring of H
- $\sigma(H)$ = minimum size of a color class in a $\chi(H)$ -coloring of H

Theorem (Burr 1981): For any connected G with $|G| \geq \sigma(H)$

$$R(G, H) \geq (|G| - 1)(\chi(H) - 1) + \sigma(H)$$

Proof:



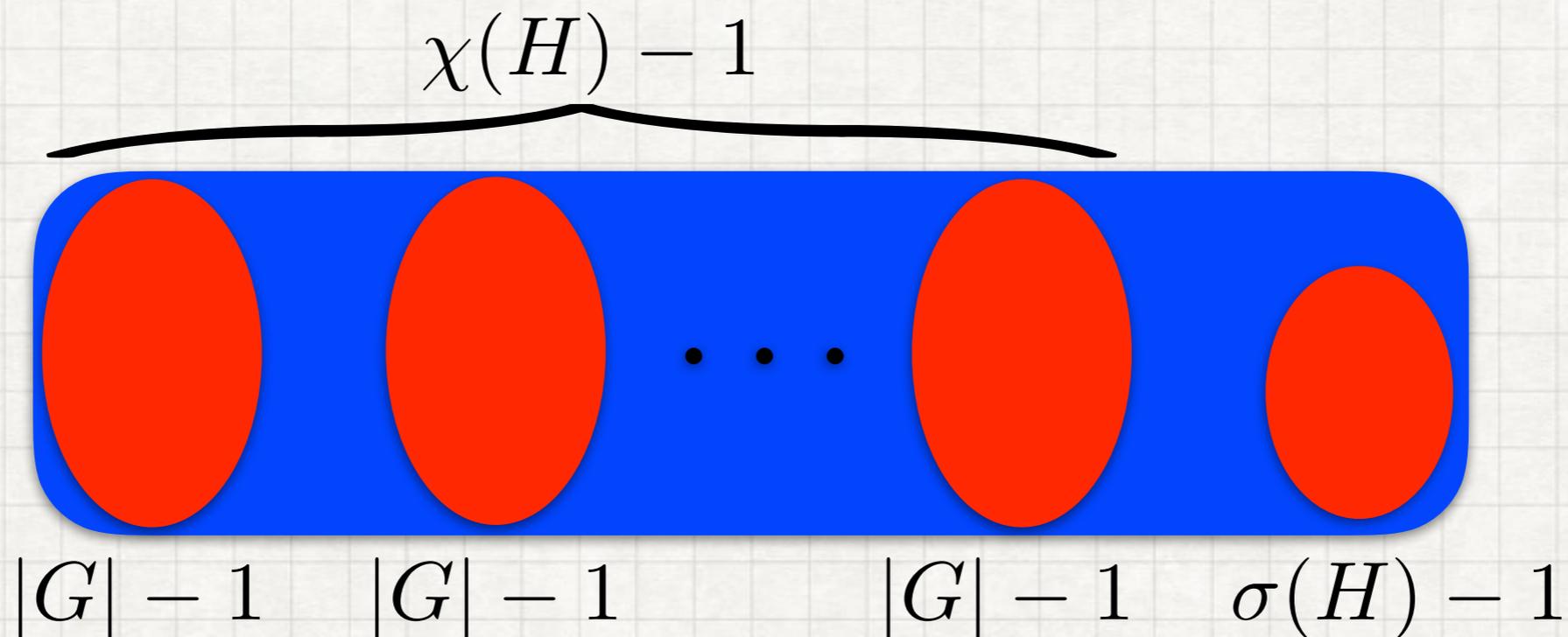
Ramsey Goodness

- $\chi(H)$ = smallest number of colors in a proper coloring of H
- $\sigma(H)$ = minimum size of a color class in a $\chi(H)$ -coloring of H

Theorem (Burr 1981): For any connected G with $|G| \geq \sigma(H)$

$$R(G, H) \geq (|G| - 1)(\chi(H) - 1) + \sigma(H)$$

Proof:



Definition: A graph G is called H -good if equality holds above.

Conjecture (Allen, Brightwell, and Skokan 2013): For $n \geq \chi(H)|H|$

$$R(P_n, H) = (n - 1)(\chi(H) - 1) + \sigma(H)$$

Conjecture (Allen, Brightwell, and Skokan 2013): For $n \geq \chi(H)|H|$

$$R(P_n, H) = (n - 1)(\chi(H) - 1) + \sigma(H)$$

- Proven when $n \geq 4|H|$ (Pokrovskiy and Sudakov 2016).

Conjecture (Allen, Brightwell, and Skokan 2013): For $n \geq \chi(H)|H|$

$$R(P_n, H) = (n - 1)(\chi(H) - 1) + \sigma(H)$$

• Proven when $n \geq 4|H|$ (Pokrovskiy and Sudakov 2016).

Theorem (Erdős, Faudree, Rousseau, Schelp 1985): For any graph H and Δ , there exists n_0 such that for all $n \geq n_0$, any tree T on n vertices with max degree Δ satisfies

$$R(T, H) = (n - 1)(\chi(H) - 1) + \sigma(H).$$

Conjecture (Allen, Brightwell, and Skokan 2013): For $n \geq \chi(H)|H|$

$$R(P_n, H) = (n - 1)(\chi(H) - 1) + \sigma(H)$$

• Proven when $n \geq 4|H|$ (Pokrovskiy and Sudakov 2016).

Theorem (Erdős, Faudree, Rousseau, Schelp 1985): For any graph H and Δ , there exists n_0 such that for all $n \geq n_0$, any tree T on n vertices with max degree Δ satisfies

$$R(T, H) = (n - 1)(\chi(H) - 1) + \sigma(H).$$

Thinking of $\Delta, \chi(H)$ as constants, their methods can at best give $n_0 = \Omega(|H|^4)$.

Conjecture (Allen, Brightwell, and Skokan 2013): For $n \geq \chi(H)|H|$

$$R(P_n, H) = (n - 1)(\chi(H) - 1) + \sigma(H)$$

• Proven when $n \geq 4|H|$ (Pokrovskiy and Sudakov 2016).

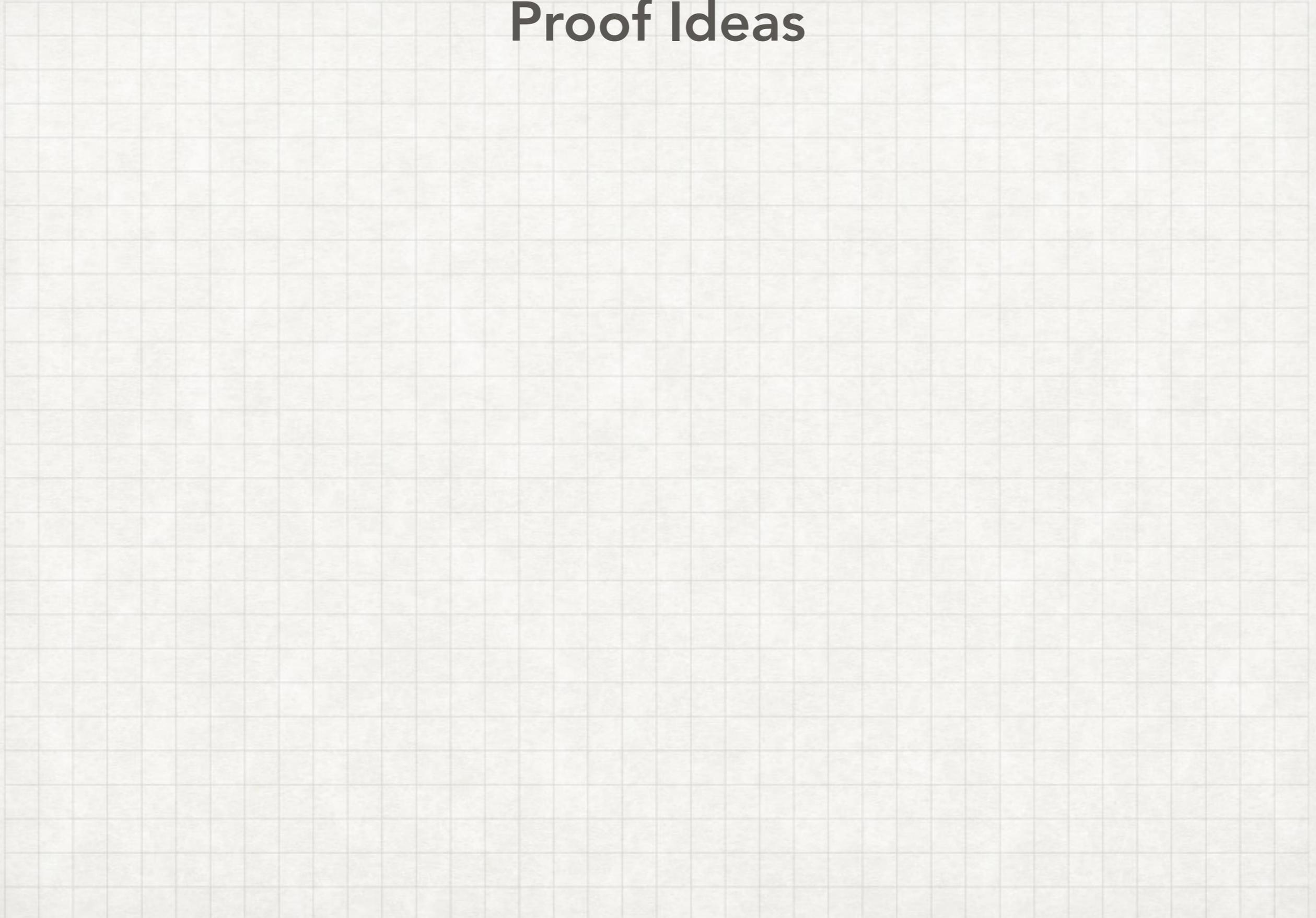
Theorem (Erdős, Faudree, Rousseau, Schelp 1985): For any graph H and Δ , there exists n_0 such that for all $n \geq n_0$, any tree T on n vertices with max degree Δ satisfies

$$R(T, H) = (n - 1)(\chi(H) - 1) + \sigma(H).$$

Thinking of $\Delta, \chi(H)$ as constants, their methods can at best give $n_0 = \Omega(|H|^4)$.

Theorem (B., Pokrovskiy, Sudakov 2016): The above theorem holds for $n_0 = \Omega(|H| \log^4 |H|)$.

Proof Ideas



Proof Ideas

For simplicity let's consider $H = K_{m,m}$ the complete bipartite graph with m vertices in each part.

Proof Ideas

For simplicity let's consider $H = K_{m,m}$ the complete bipartite graph with m vertices in each part.

Q: If we have a **red-blue** complete graph on $n - 1 + m$ vertices and the **blue** graph has no copy of $K_{m,m}$, what can we say about the **red** graph?

Proof Ideas

For simplicity let's consider $H = K_{m,m}$ the complete bipartite graph with m vertices in each part.

Q: If we have a **red-blue** complete graph on $n - 1 + m$ vertices and the **blue** graph has no copy of $K_{m,m}$, what can we say about the **red** graph?

A: The **red** graph is an expander! (almost)

Proof Ideas

For simplicity let's consider $H = K_{m,m}$ the complete bipartite graph with m vertices in each part.

Q: If we have a **red-blue** complete graph on $n - 1 + m$ vertices and the **blue** graph has no copy of $K_{m,m}$, what can we say about the **red** graph?

A: The **red** graph is an expander! (almost)

For any set S of m vertices, $|N(S)| \geq n - m$.

Proof Ideas

For simplicity let's consider $H = K_{m,m}$ the complete bipartite graph with m vertices in each part.

Q: If we have a **red-blue** complete graph on $n - 1 + m$ vertices and the **blue** graph has no copy of $K_{m,m}$, what can we say about the **red** graph?

A: The **red** graph is an expander! (almost)

For any set S of m vertices, $|N(S)| \geq n - m$.

Theorem (Haxell 2001): In an expander on n vertices, we can find any bounded degree tree on $.99n$ vertices.

Proof Ideas

For simplicity let's consider $H = K_{m,m}$ the complete bipartite graph with m vertices in each part.

Q: If we have a **red-blue** complete graph on $n - 1 + m$ vertices and the **blue** graph has no copy of $K_{m,m}$, what can we say about the **red** graph?

A: The **red** graph is an expander! (almost)

For any set S of m vertices, $|N(S)| \geq n - m$.

Theorem (Haxell 2001): In an expander on n vertices, we can find any bounded degree tree on $.99n$ vertices.

Theorem (Montgomery 2014): For any tree T on n vertices with max degree Δ , the random graph $G(n, \Delta \log^5 n/n)$ almost surely contains a copy of T .

Conjecture: For any tree T on n vertices with max degree Δ ,
and any graph H with $n \geq O(\Delta|H|)$,

$$R(T, H) = (n - 1)(\chi(H) - 1) + \sigma(H).$$

Conjecture: For any tree T on n vertices with max degree Δ ,
and any graph H with $n \geq O(\Delta|H|)$,

$$R(T, H) = (n - 1)(\chi(H) - 1) + \sigma(H).$$

- We can prove it for trees with linearly many leaves.

Conjecture: For any tree T on n vertices with max degree Δ ,
and any graph H with $n \geq O(\Delta|H|)$,

$$R(T, H) = (n - 1)(\chi(H) - 1) + \sigma(H).$$

- We can prove it for trees with linearly many leaves.

Thank you.